

A tour of String Diagrams and Monoidal Categories

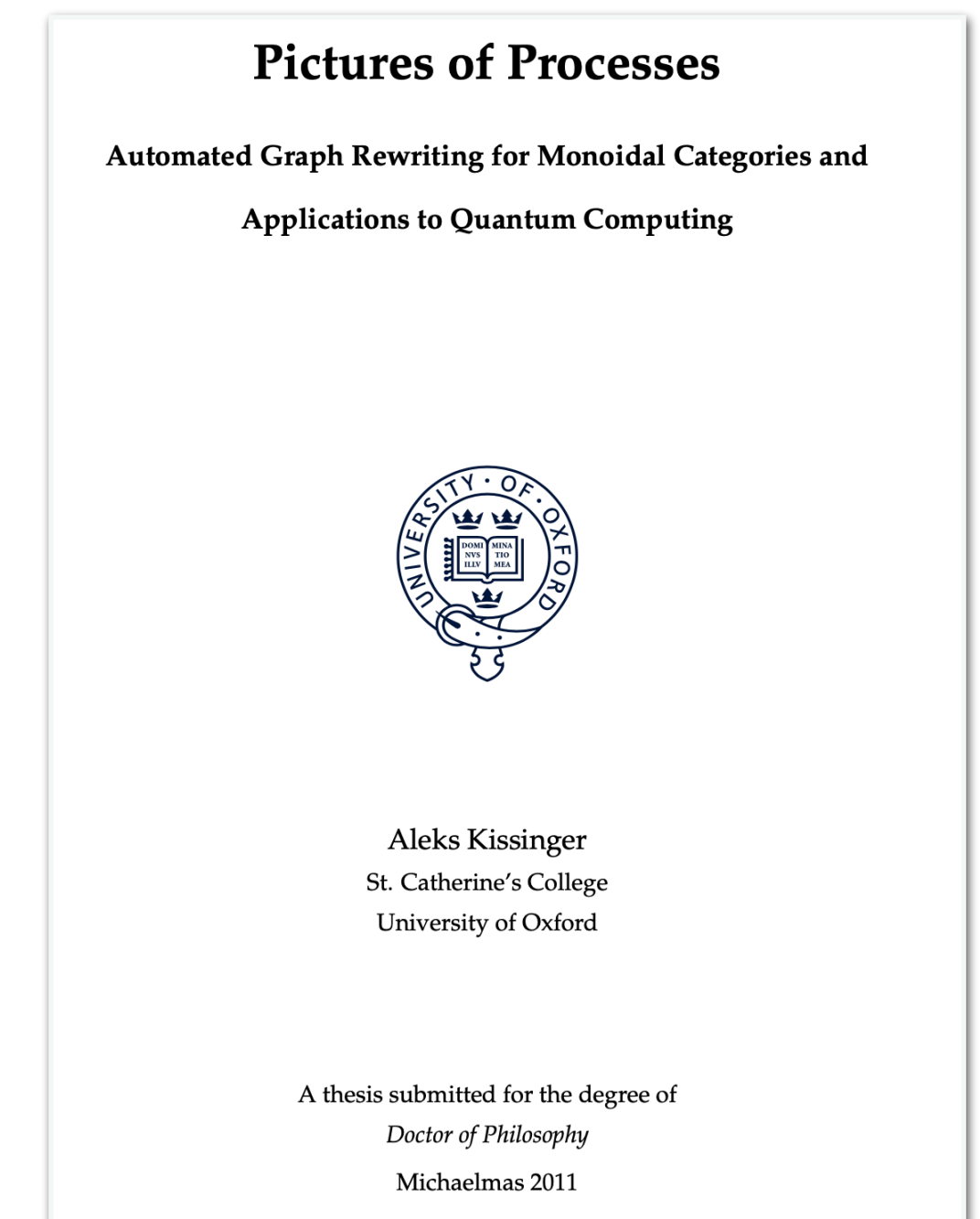
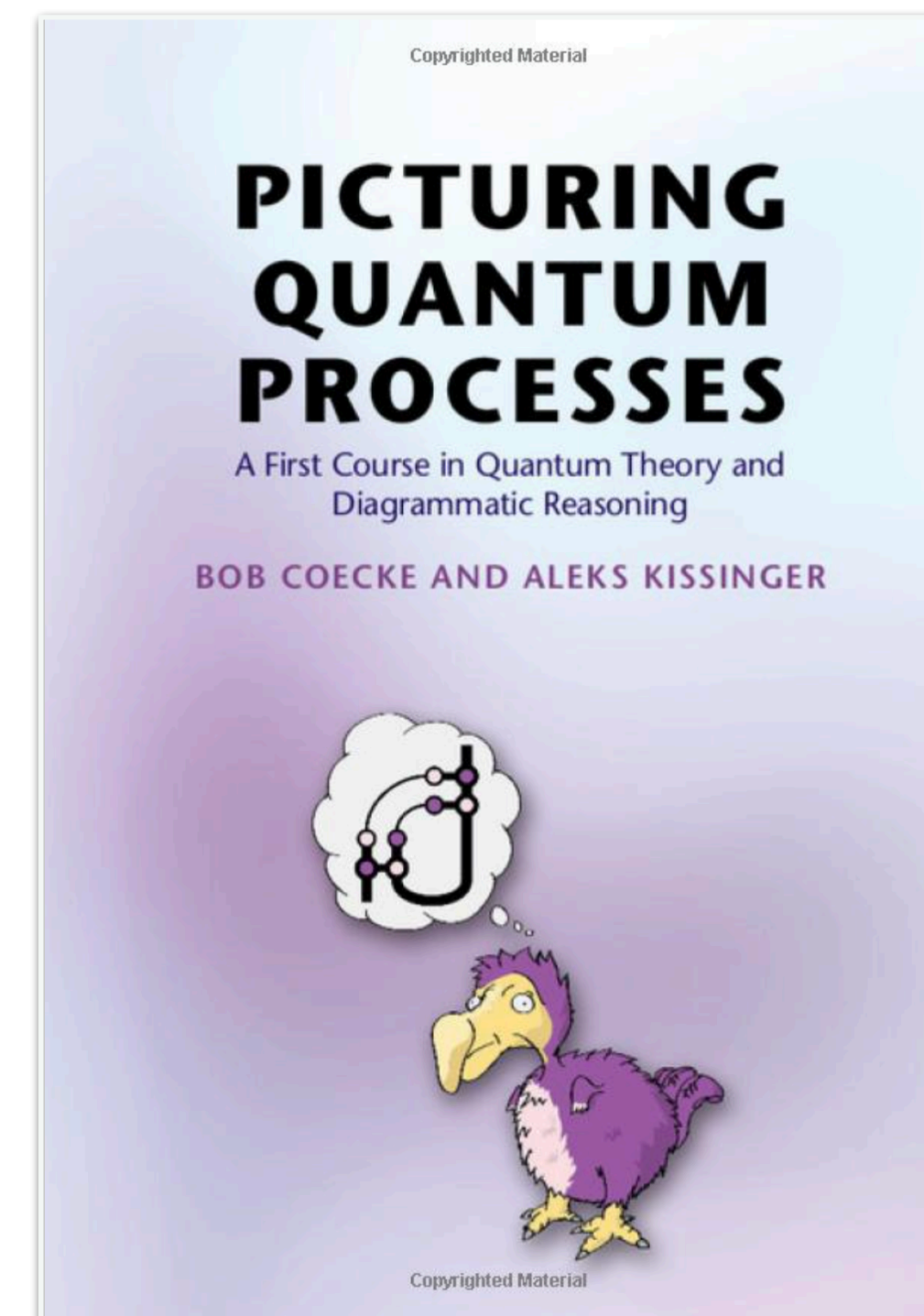
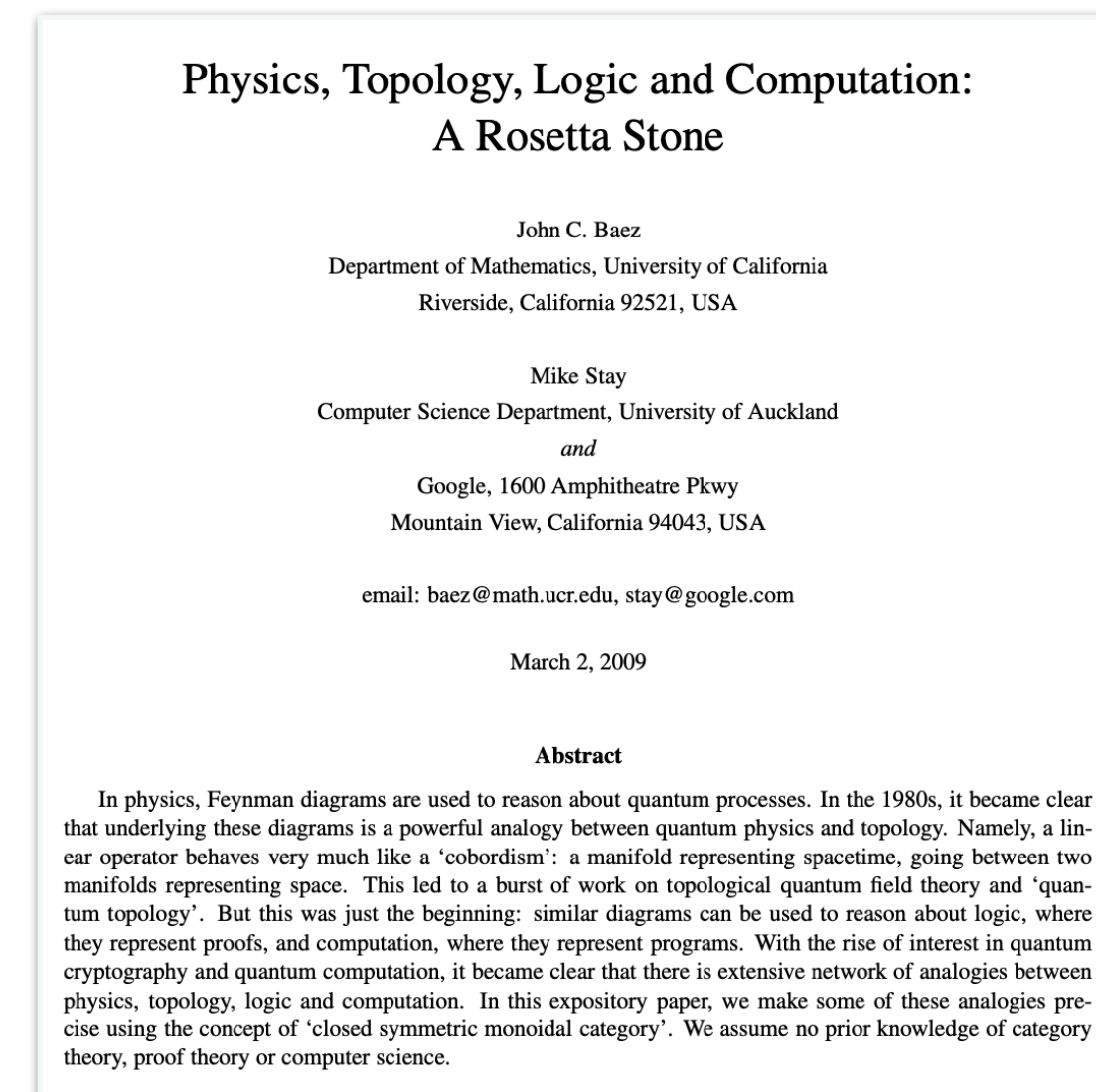
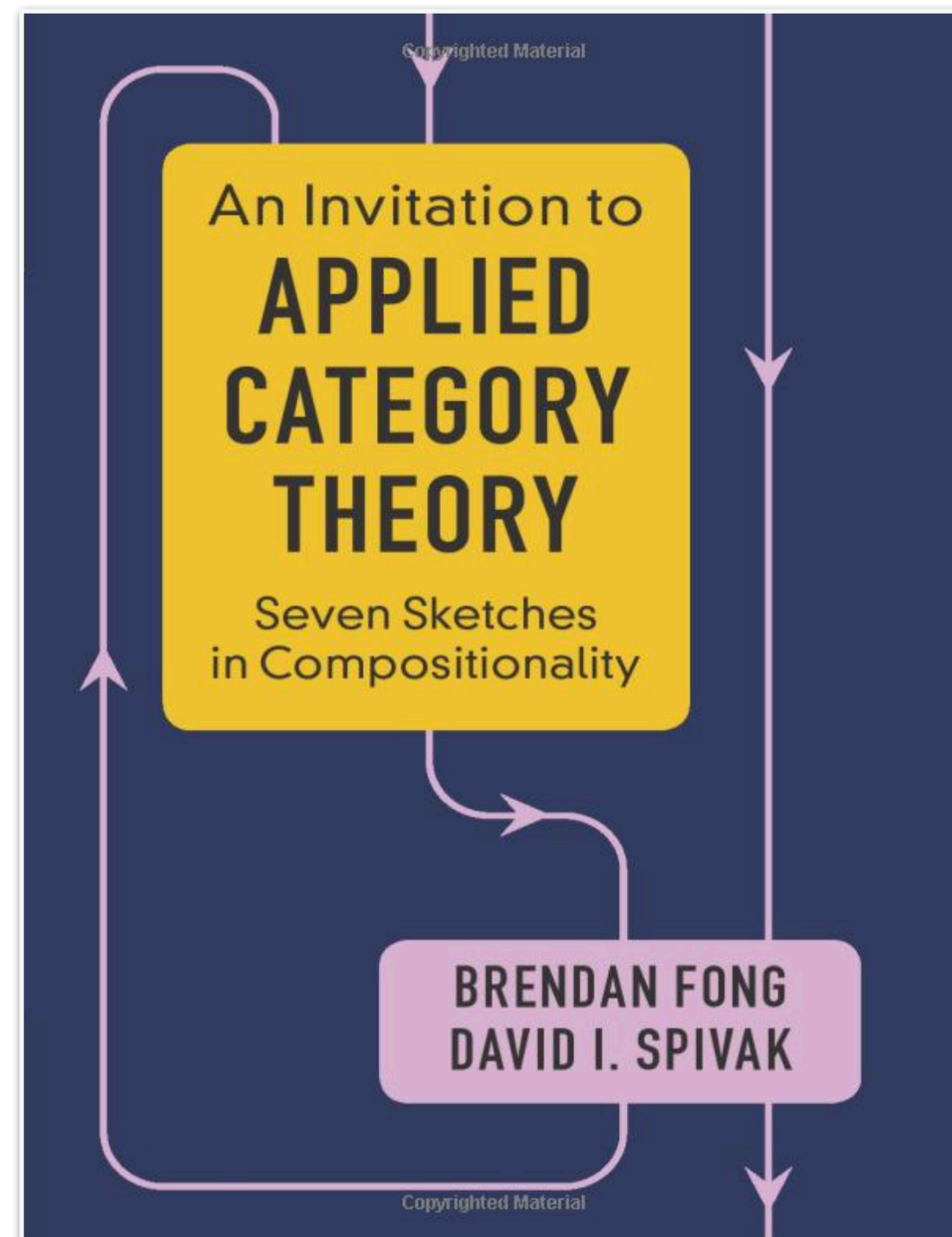
Juan Pablo Romero Méndez
@1jpablo1

Scale by the Bay
November 2020

Goals

- Show why Monoidal Categories are a good modeling tool for certain class of domains
- Show how Monoidal Categories can be represented graphically via String Diagrams
- Given the time constraint I'll go rather quickly but don't hesitate to ask questions after the presentation!

A few references



Part I:
String Diagrams as
Processes

Processes

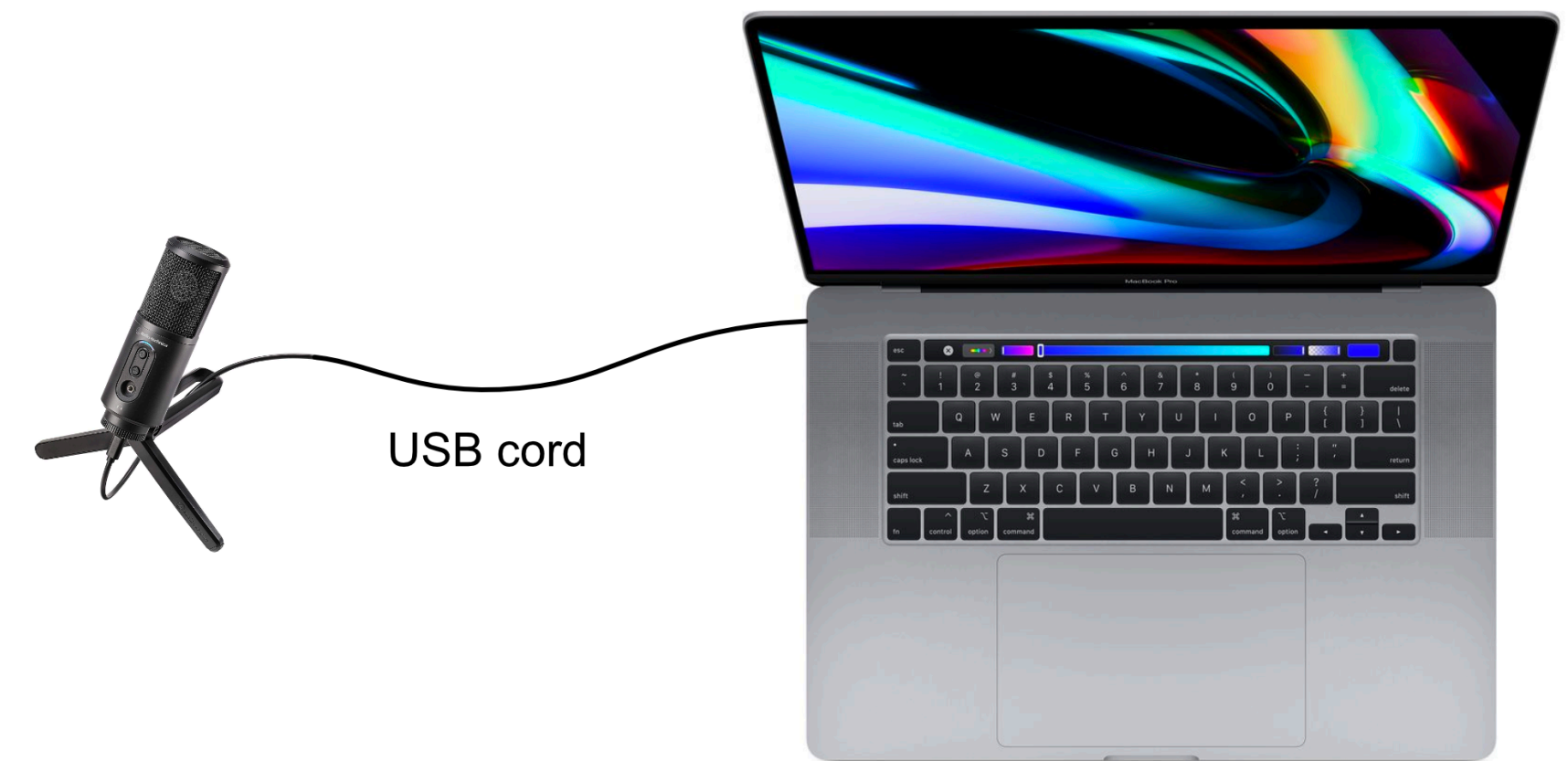
sort: List[A] \Rightarrow List[A]

$$M = \begin{pmatrix} a & b & c \\ d & e & f \end{pmatrix}$$

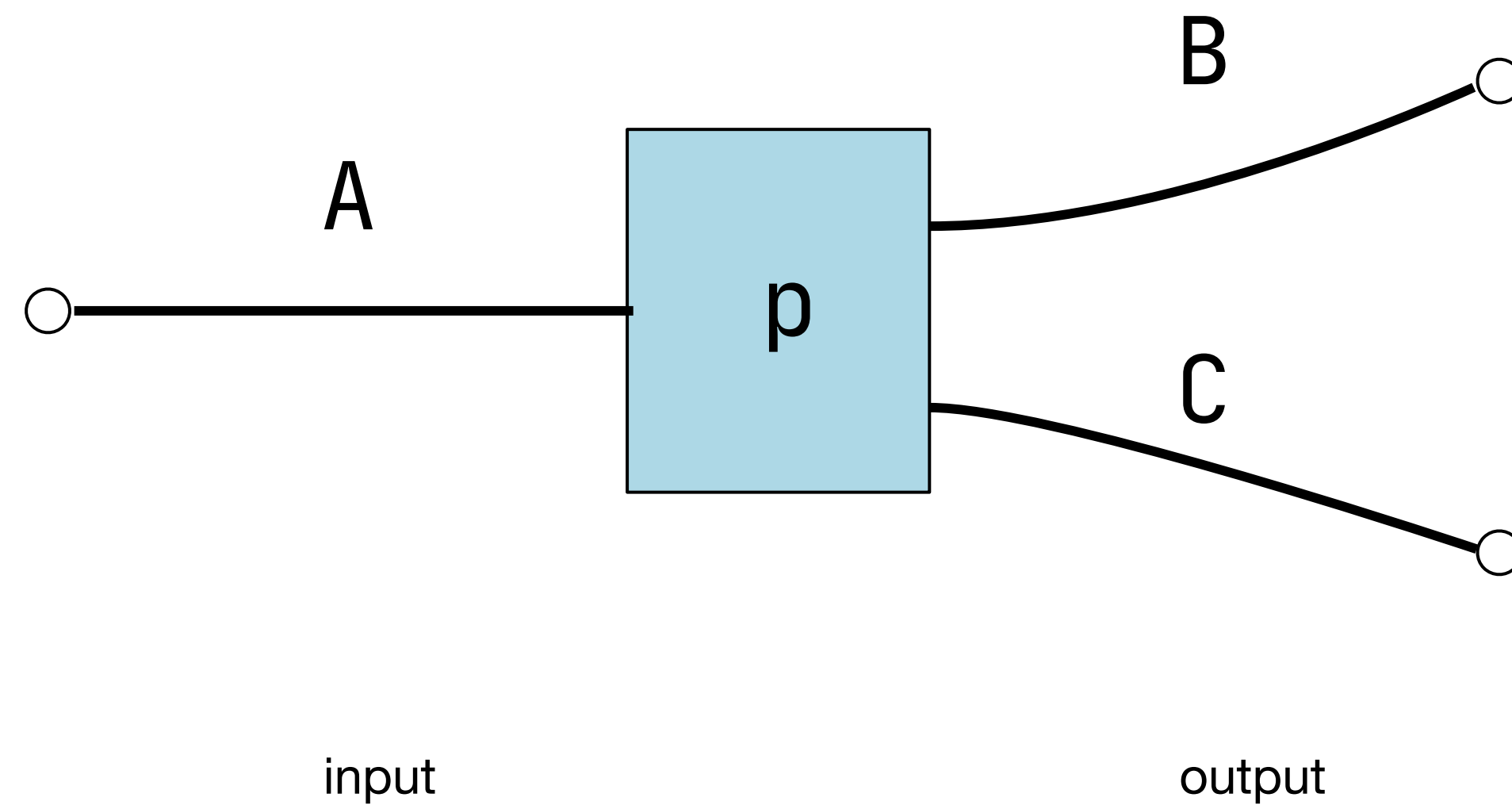
$$f(x, y) = x^2 + y^2$$

INGREDIENTS

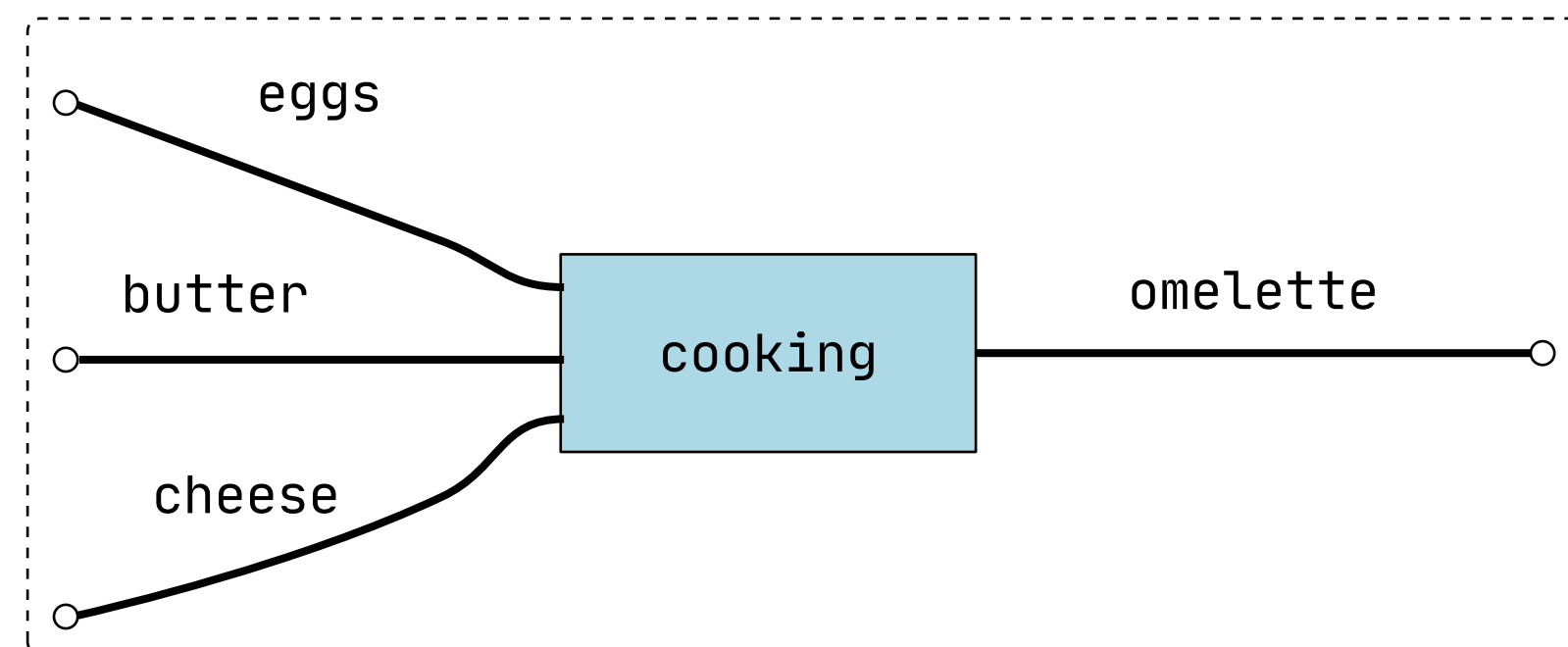
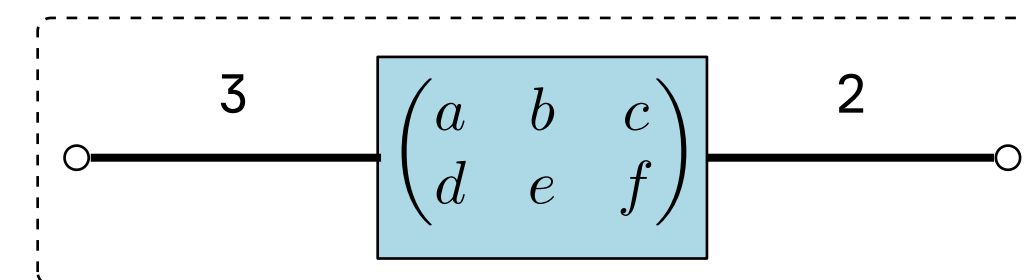
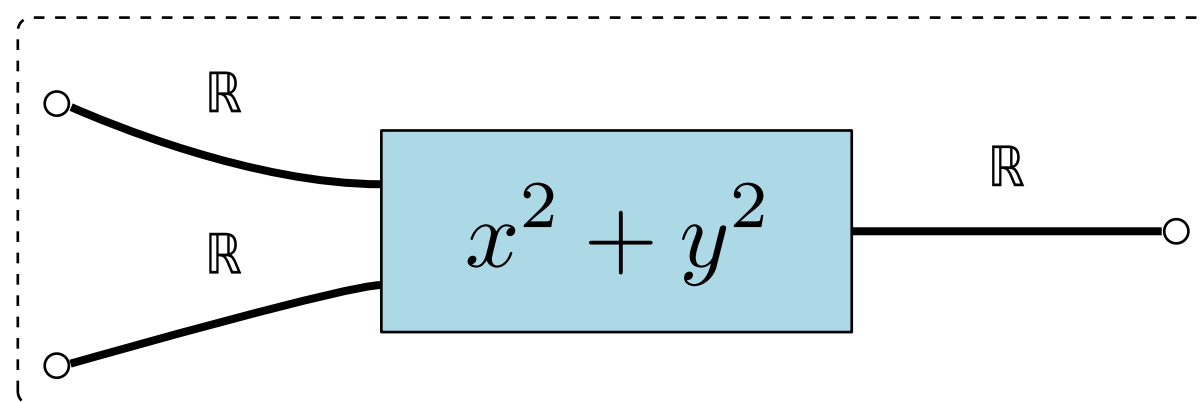
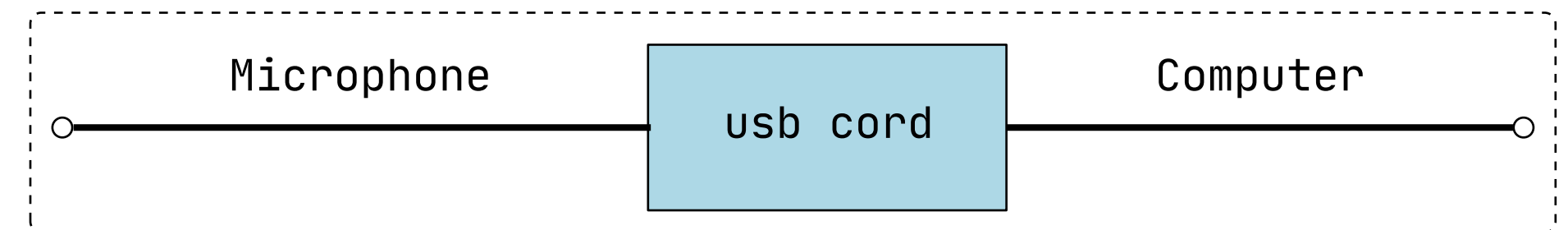
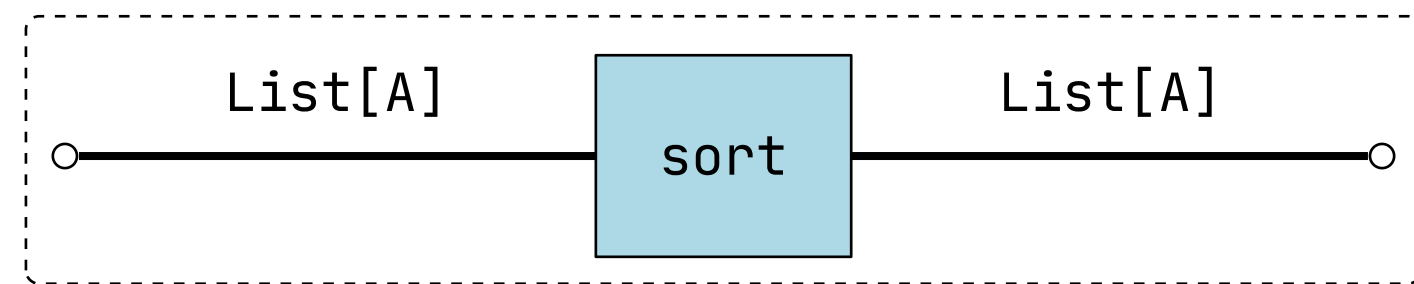
2 large eggs
Kosher salt
Freshly ground black pepper
Pinch red pepper flakes
2 tbsp. butter
1/4 c. shredded cheddar
2 tbsp. freshly chopped chives



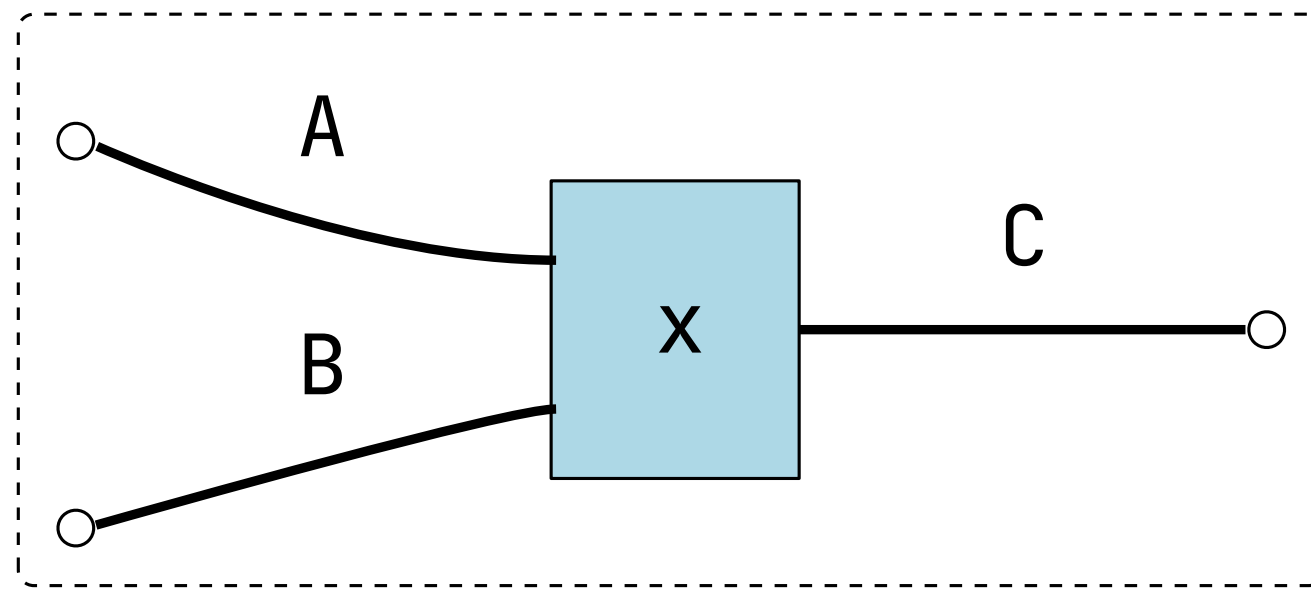
A process



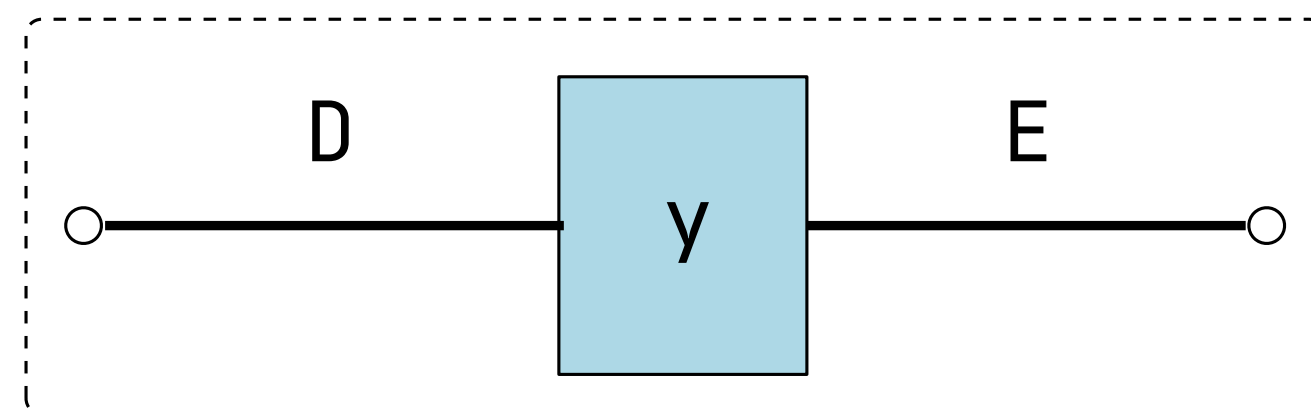
Processes as diagrams



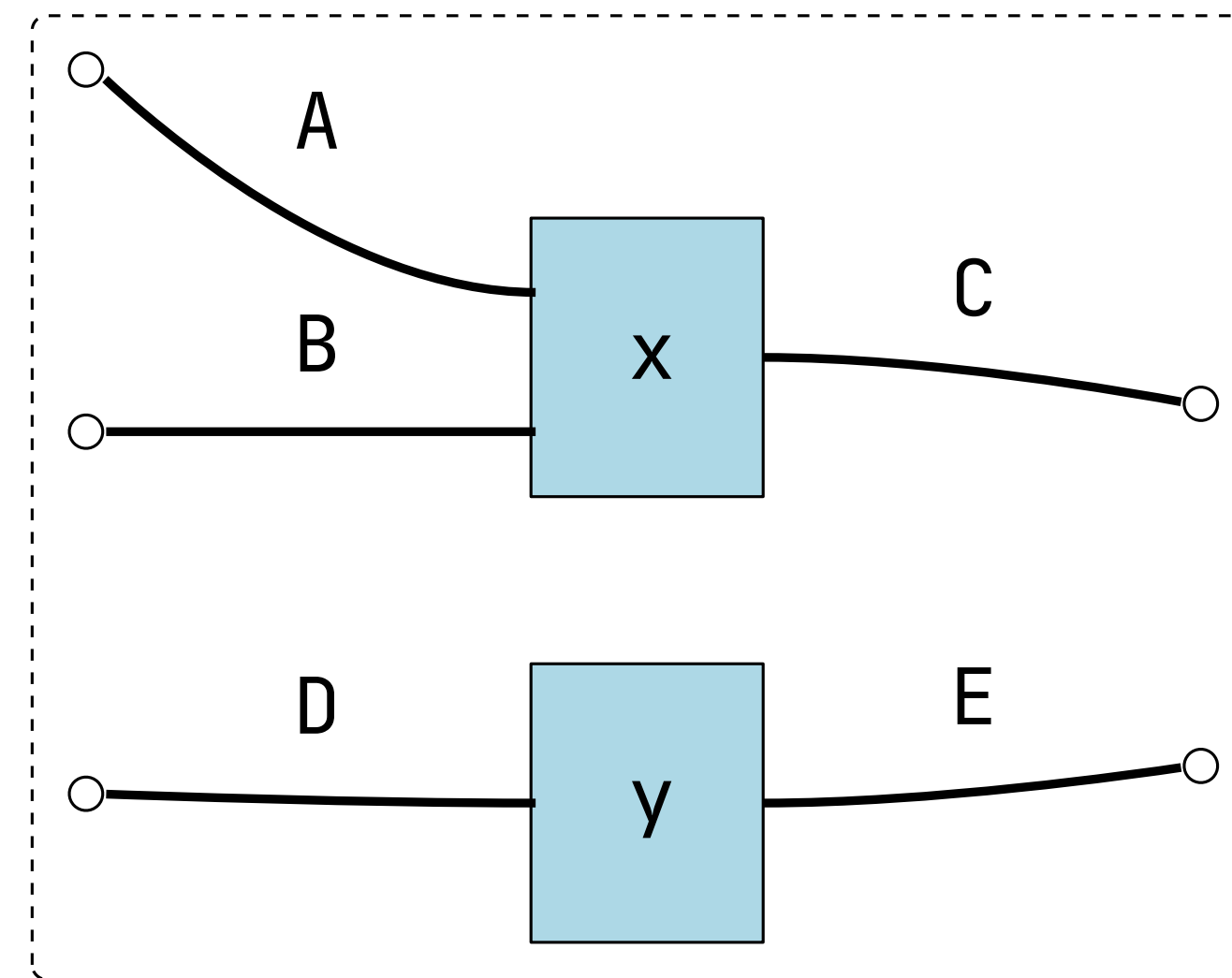
Parallel composition



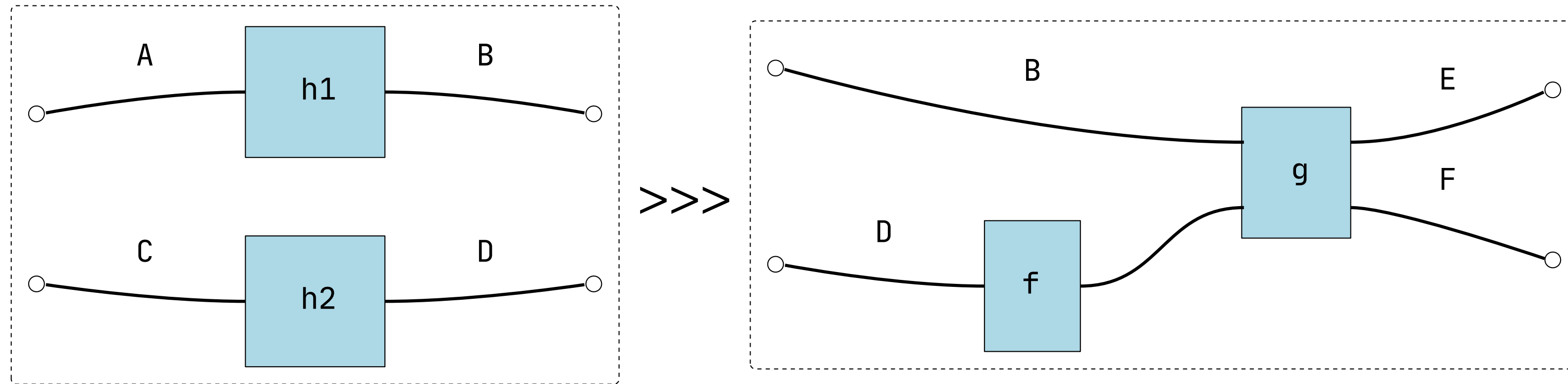
\otimes



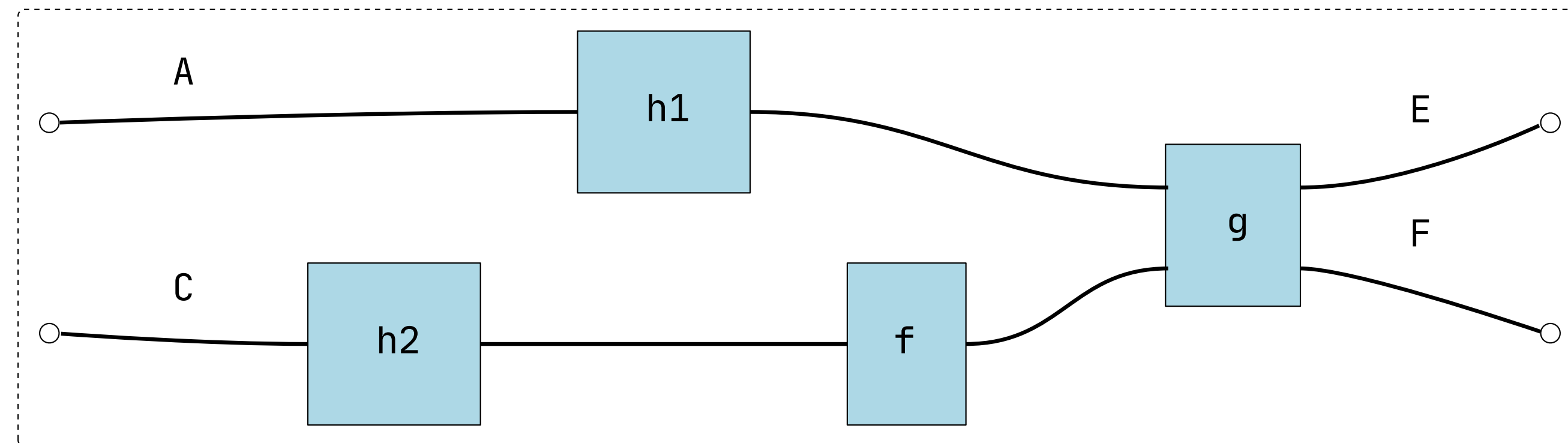
=



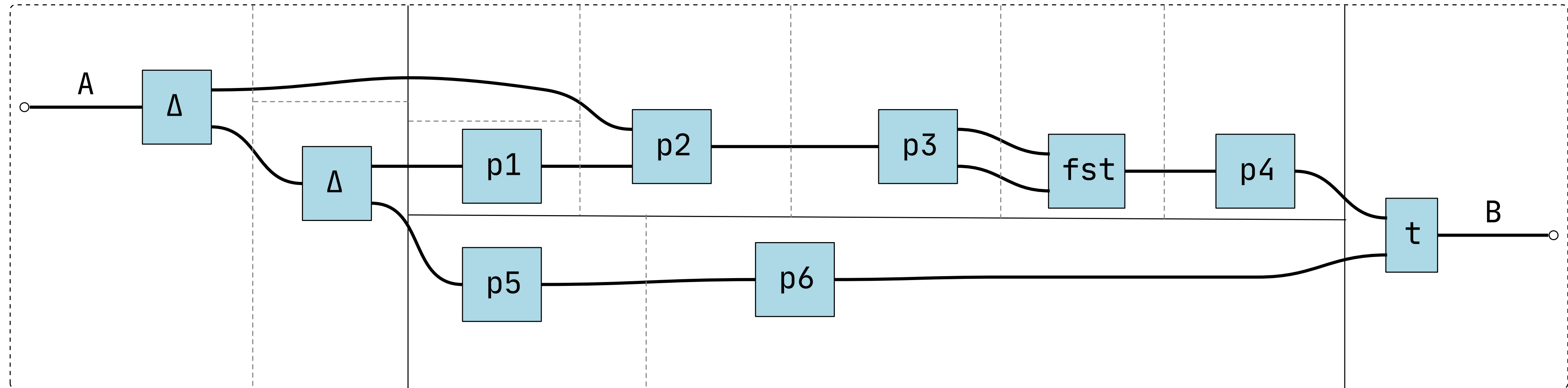
Sequential composition



=



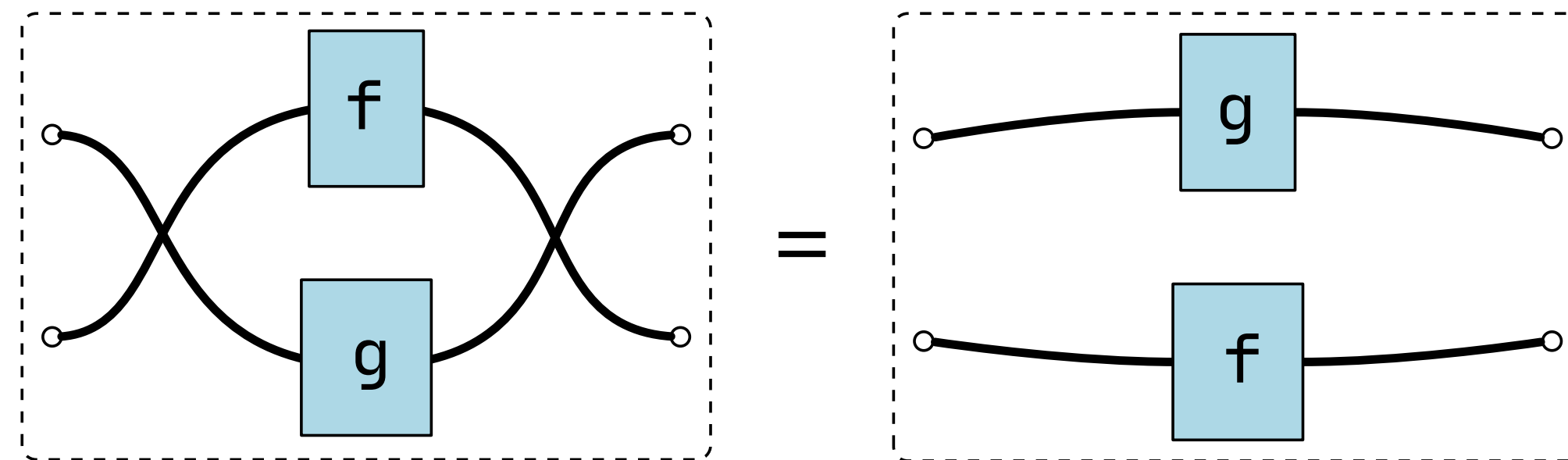
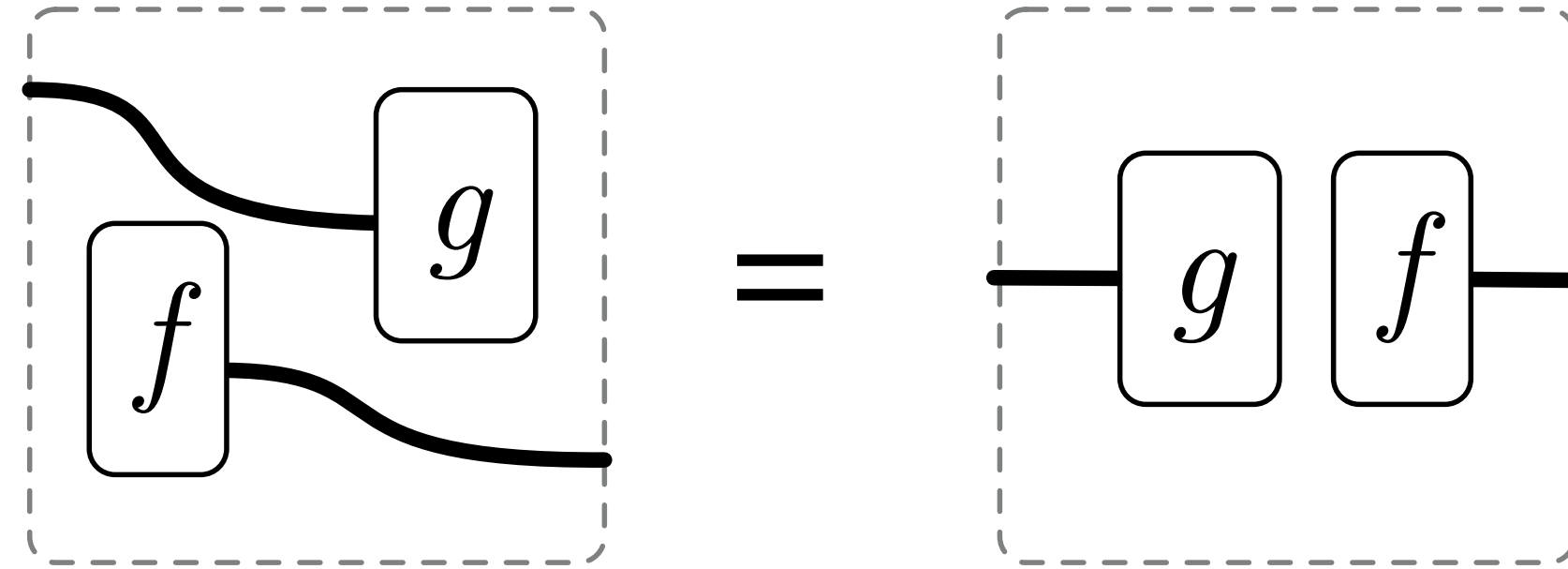
Complex diagrams



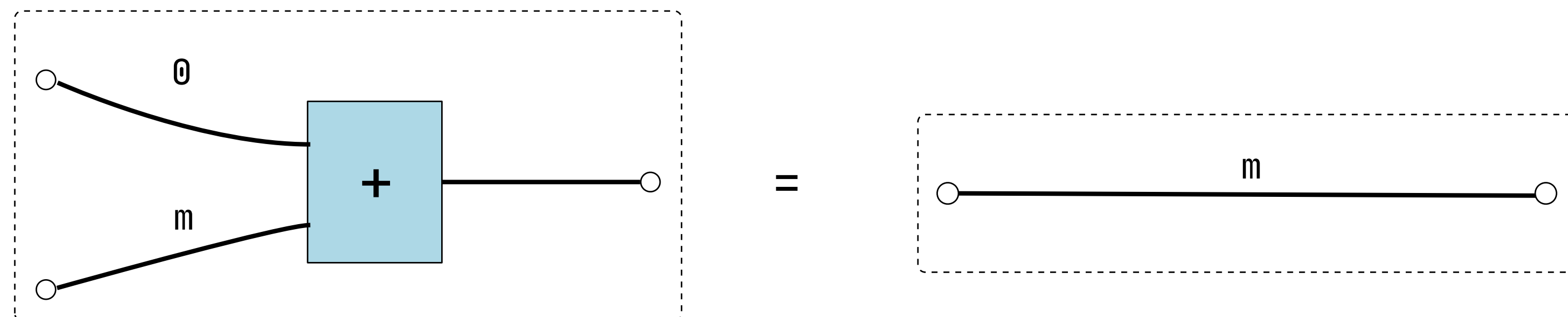
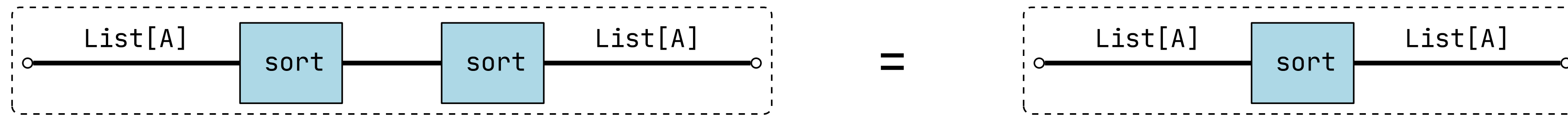
Process theories

- A "process theory" is an interpretation of a diagram in terms of a concrete class of processes.
- In particular it provides an interpretation of **wires**, **boxes**, and **composition** operations of diagrams.
- Examples of process theories:
 - Functions and sets
 - linear maps and vector spaces
 - Matrices and natural numbers

Diagram equations



Process equations

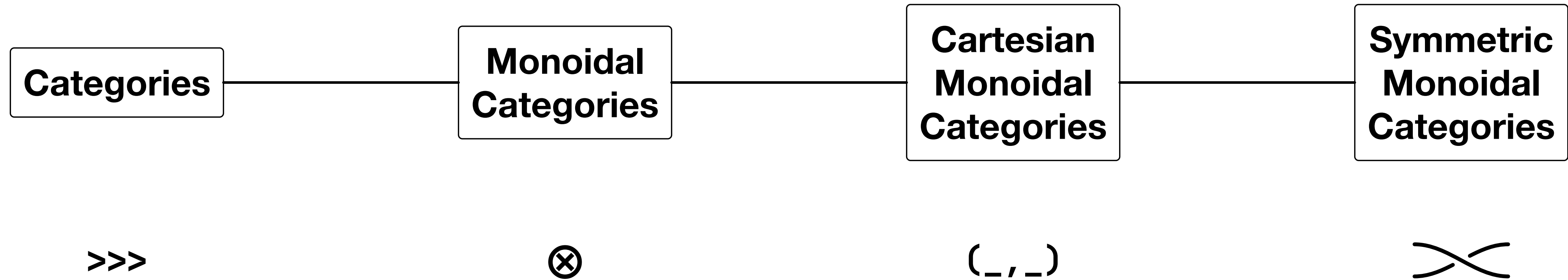


Part II:
String Diagrams as
Monoidal Categories

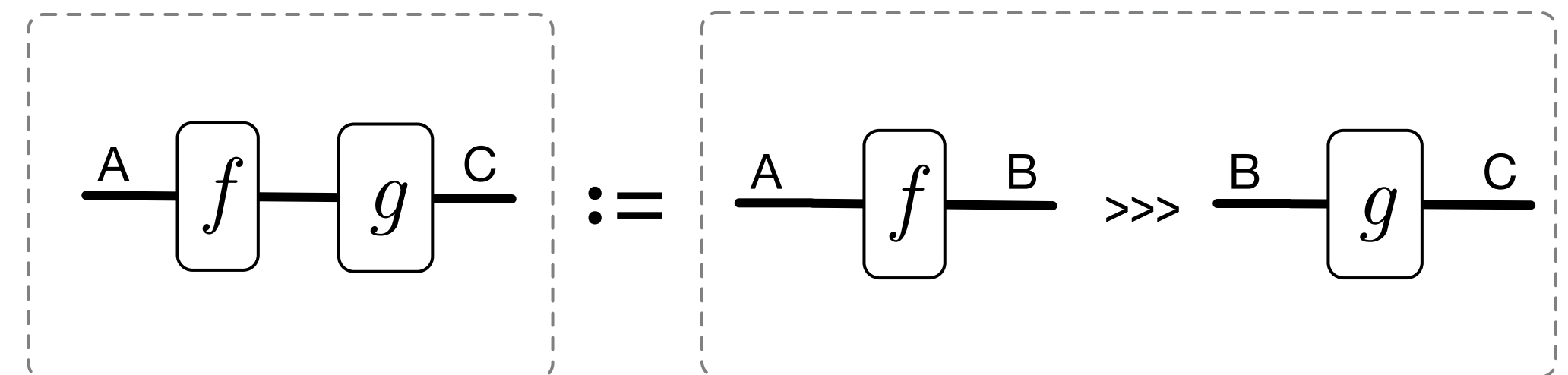
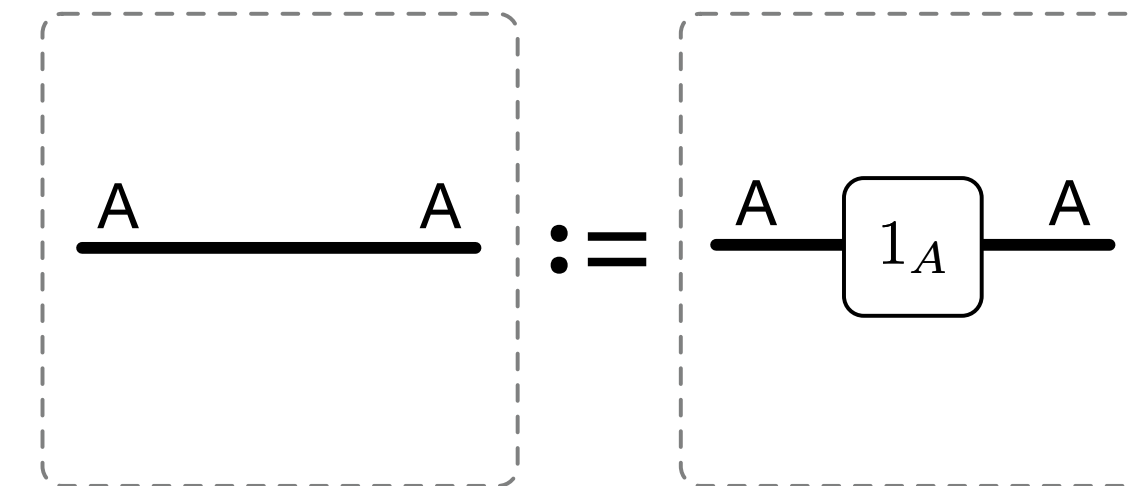
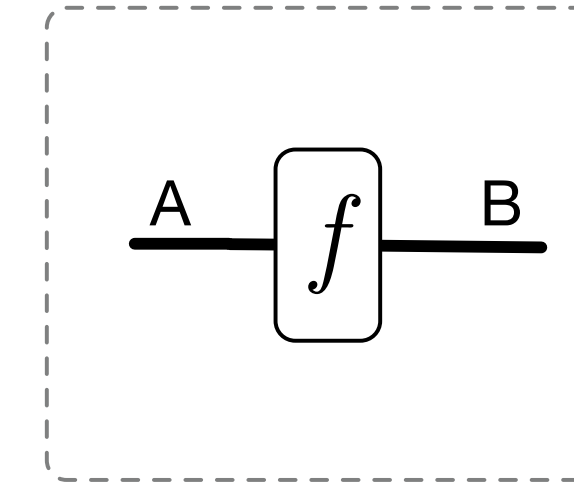
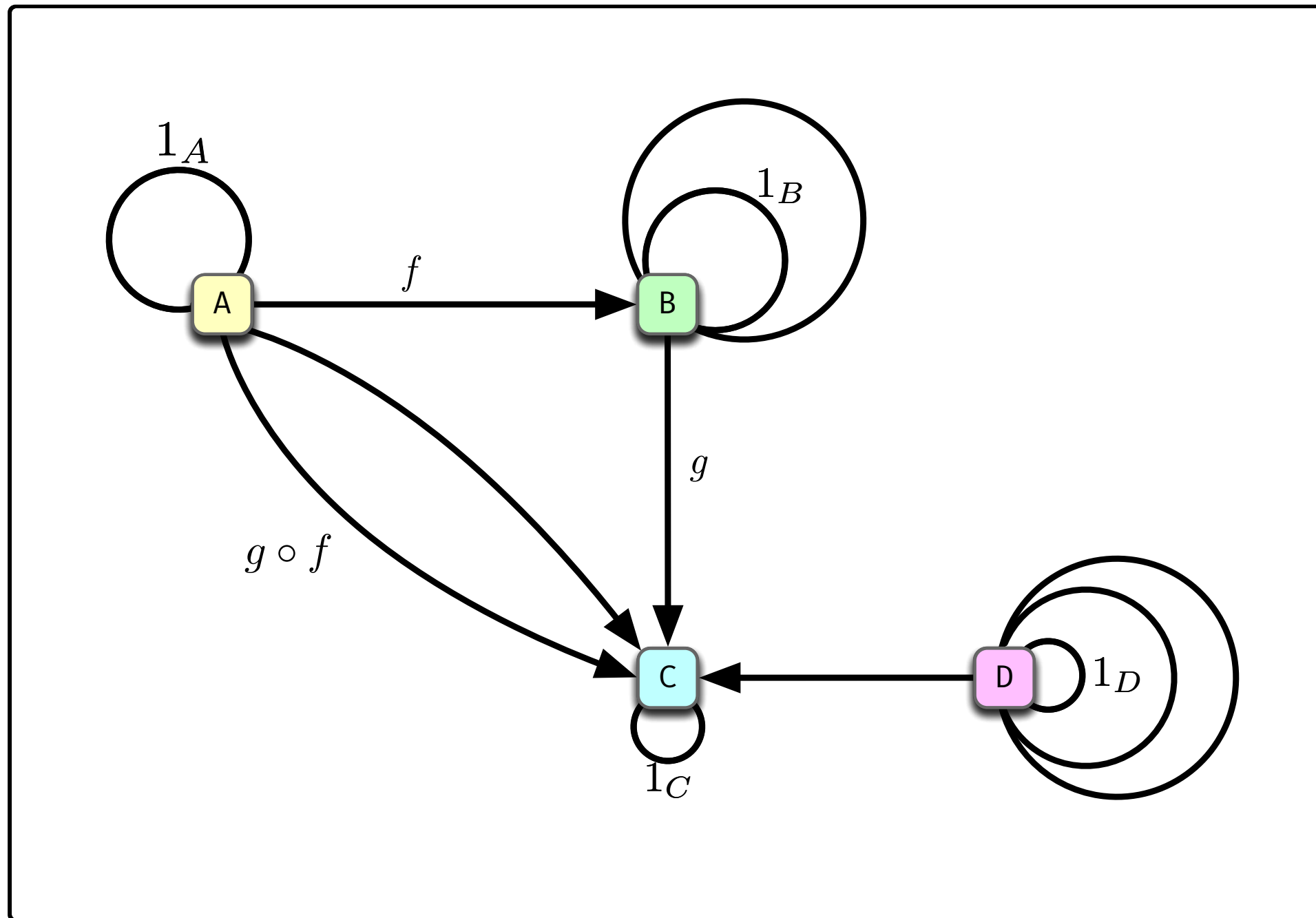
Formalization

This diagrammatic language can be formalized by using category theory, specifically **monoidal categories**.

Plan



1. Categories



Categories: equations

$$\begin{array}{c} A \quad A \\ \hline \end{array} \ggg \begin{array}{c} A \quad \boxed{f} \quad B \\ \hline \end{array}$$

=

$$\begin{array}{c} A \quad \boxed{f} \quad B \\ \hline \end{array}$$

=

$$\begin{array}{c} A \quad \boxed{f} \quad B \\ \hline \end{array} \ggg \begin{array}{c} B \quad B \\ \hline \end{array}$$

$$\left(\begin{array}{c} A \quad \boxed{f} \quad B \\ \hline \end{array} \ggg \begin{array}{c} B \quad \boxed{g} \quad C \\ \hline \end{array} \right) \ggg \begin{array}{c} C \quad \boxed{h} \quad D \\ \hline \end{array}$$

=

$$\begin{array}{c} A \quad \boxed{f} \quad \boxed{g} \quad \boxed{h} \quad D \\ \hline \end{array}$$

=

$$\begin{array}{c} A \quad \boxed{f} \quad B \\ \hline \end{array} \ggg \left(\begin{array}{c} B \quad \boxed{g} \quad C \\ \hline \end{array} \ggg \begin{array}{c} C \quad \boxed{h} \quad D \\ \hline \end{array} \right)$$

Example: Matrices

- objects: natural numbers
- arrows $n \rightarrow m$:
matrices $m \times n$
- composition:
matrix multiplication
- identity 1_n :
identity matrix I_n

$$\begin{array}{c}
 \text{input } (n) \downarrow \\
 A = \begin{pmatrix} a_{11} & \cdots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{m1} & \cdots & a_{mn} \end{pmatrix} \xrightarrow{\text{output } (m)} \\
 \begin{array}{c} n \quad \boxed{A} \quad m \end{array}
 \end{array}$$

$$\begin{array}{c}
 B = \begin{pmatrix} b_{11} & \cdots & b_{1p} \\ \vdots & \ddots & \vdots \\ b_{n1} & \cdots & b_{np} \end{pmatrix} \\
 \begin{array}{c} p \quad \boxed{B} \quad n \end{array}
 \end{array}$$

$$\begin{array}{c} p \quad \boxed{B} \quad \boxed{A} \quad m \end{array}$$

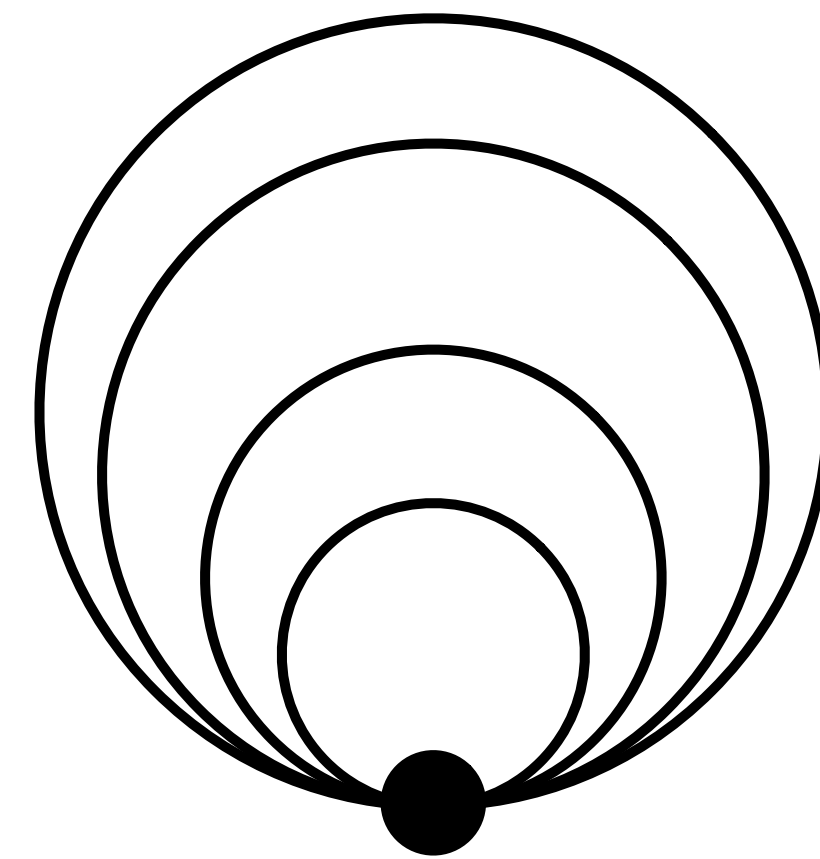
$$A \cdot B = \begin{pmatrix} c_{11} & \cdots & c_{1p} \\ \vdots & \ddots & \vdots \\ c_{m1} & \cdots & c_{mp} \end{pmatrix}$$

$$\begin{array}{c} n \quad \boxed{I} \quad n \end{array}$$

$$I_n = \begin{pmatrix} 1 & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & 1 \end{pmatrix}$$

2a. Strict Monoidal Categories

A category with a monoid structure on the objects and arrows

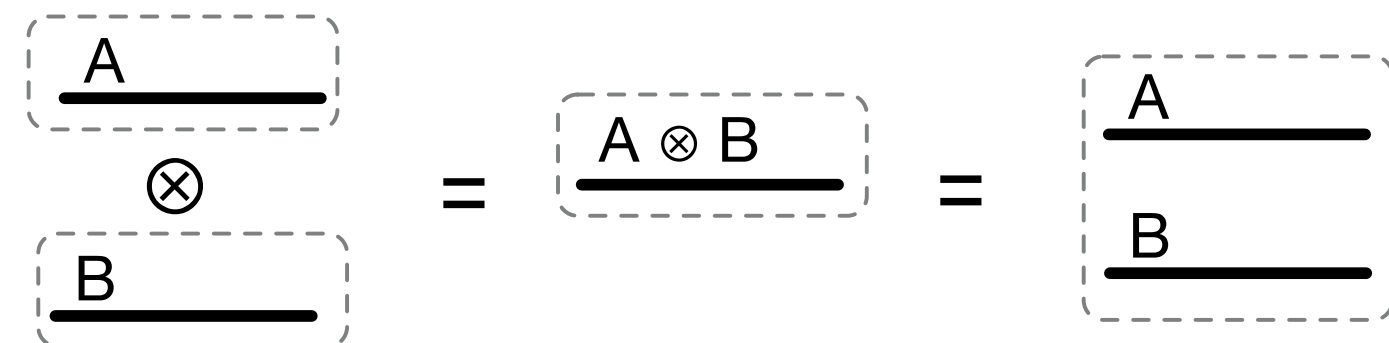


Strict Monoidal Categories

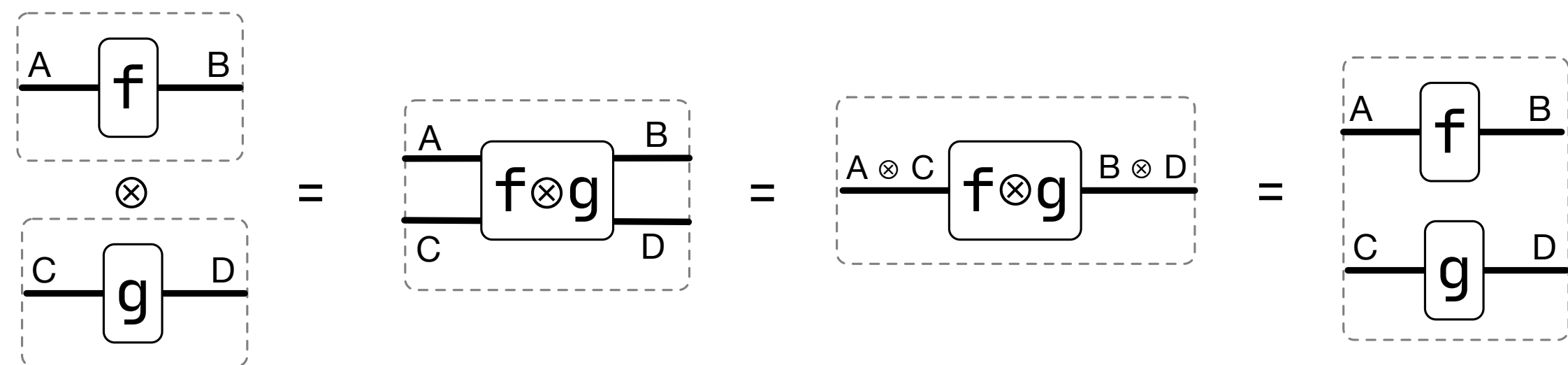
- Parallel composition: an *associative* binary operation called the **tensor** product

$$\otimes : C \times C \rightarrow C$$

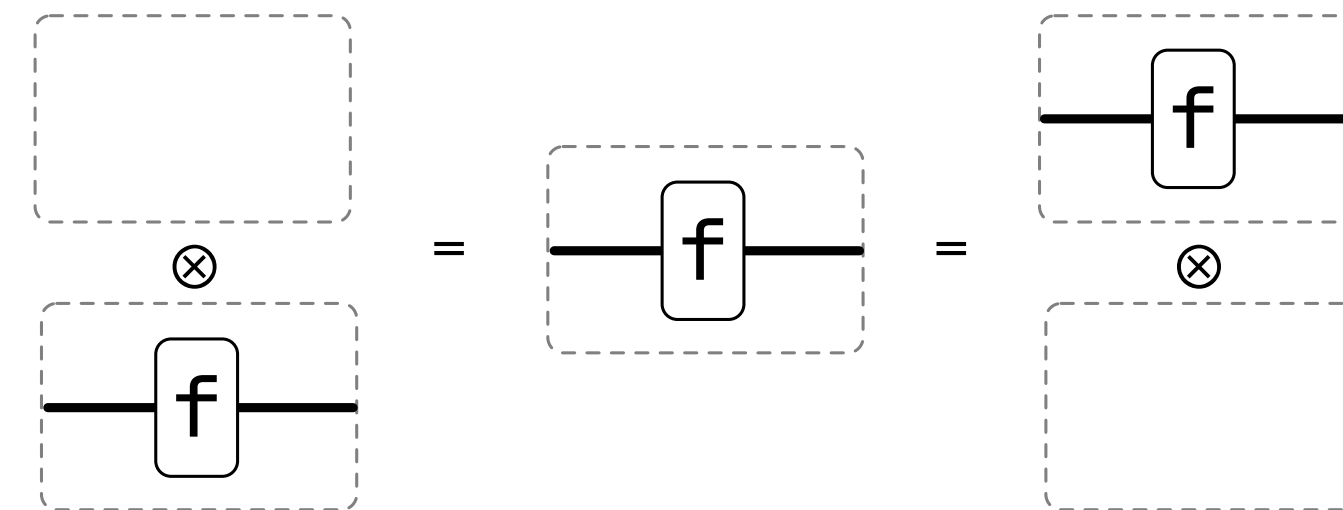
- That operates on objects



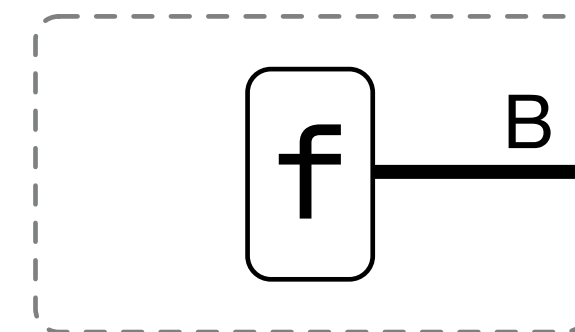
- and on arrows



- A neutral or unit object I , with identity $1_I : I \rightarrow I$

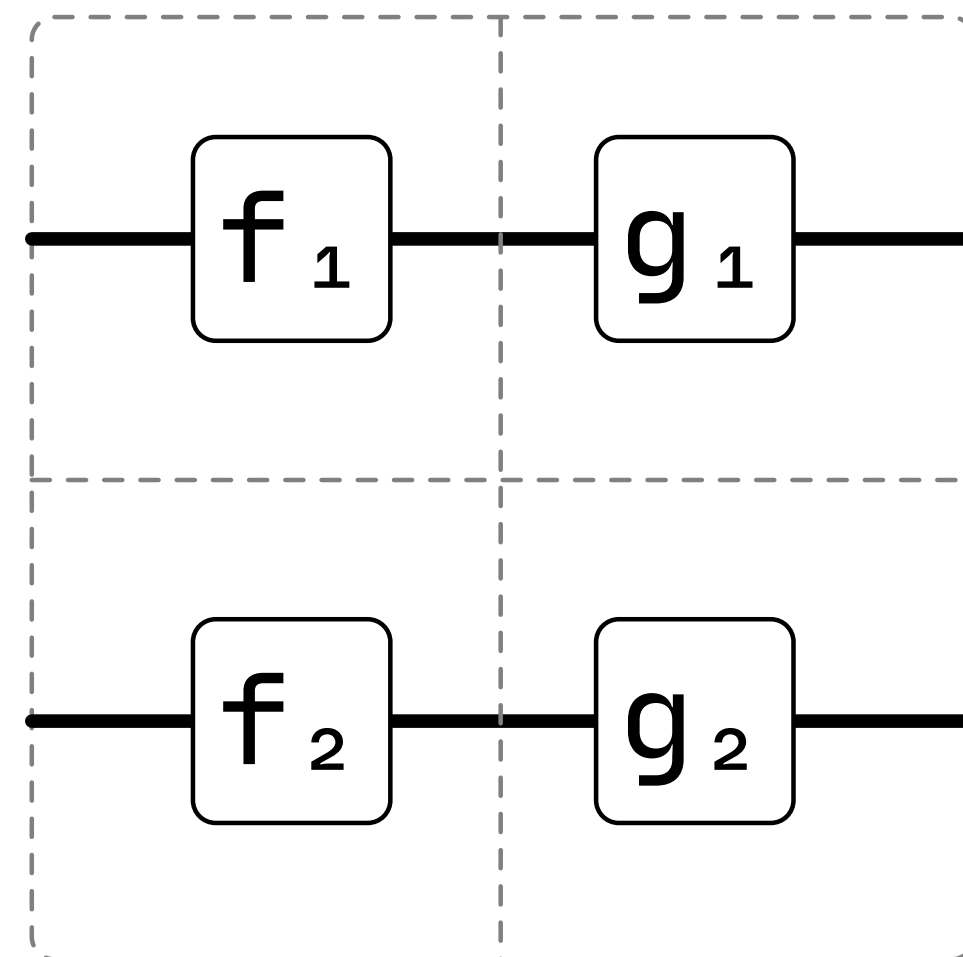


- for example, a morphism $f : I \rightarrow B$



One more law

$$(f_1 \ggg g_1) \otimes (f_2 \ggg g_2) = (f_1 \otimes f_2) \ggg (g_1 \otimes g_2)$$



Matrices again

- Matrices form a monoidal category if we use the Kronecker product as parallel composition of arrows

$$\overset{n}{\text{---}} \boxed{A} \text{---}^m$$

$$A = \begin{pmatrix} a_{11} & \cdots & a_{1n} \\ \vdots & \ddots & \vdots \\ a_{m1} & \cdots & a_{mn} \end{pmatrix}$$

$$\overset{q}{\text{---}} \boxed{B} \text{---}^p$$

$$B = \begin{pmatrix} b_{11} & \cdots & b_{1p} \\ \vdots & \ddots & \vdots \\ b_{p1} & \cdots & b_{qp} \end{pmatrix}$$

- On objects it operates as multiplication of natural numbers

$$\overset{qn}{\text{---}} \boxed{A \otimes B} \text{---}^{pm}$$

$$A \otimes B = \begin{pmatrix} a_{11}B & \cdots & a_{1n}B \\ \vdots & \ddots & \vdots \\ a_{m1}B & \cdots & a_{mn}B \end{pmatrix}$$

$$I = ()$$

- The number 0 is the unit I

2b. Monoidal categories

- In many cases it's not literally true that:

$$(X \otimes Y) \otimes Z = X \otimes (Y \otimes Z)$$

or

$$I \otimes X = X = X \otimes I$$

- For example in Scala

$$((x, y), z) \neq (x, (y, z))$$

- Instead we're forced to require isomorphisms:

- **associator:**

$$a_{X,Y,Z} : (X \otimes Y) \otimes Z \xrightarrow{\sim} X \otimes (Y \otimes Z)$$

- left and right **unitors:**

$$l_X : I \otimes X \xrightarrow{\sim} X$$

$$r_X : X \otimes I \xrightarrow{\sim} X$$

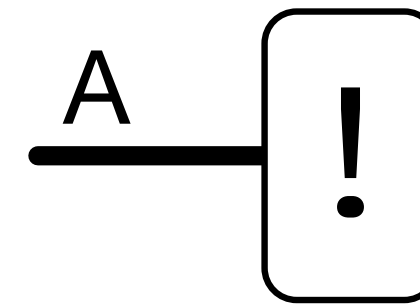
- *(satisfying certain laws)*

3. Cartesian Monoidal Categories

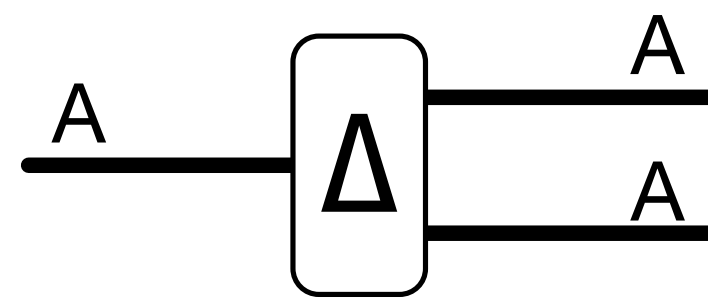
- We can extract the first and second components:



- and we can discard information:

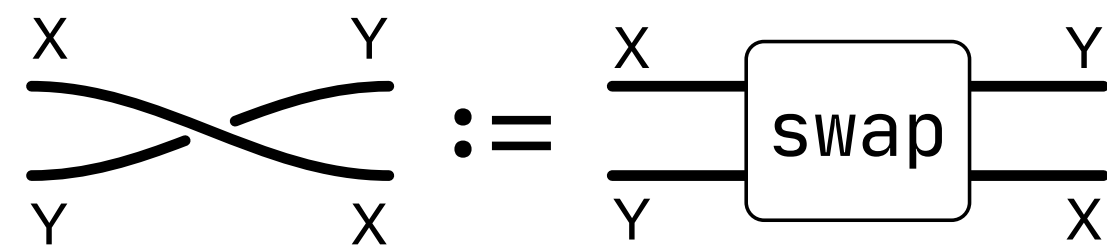


- we can duplicate information:

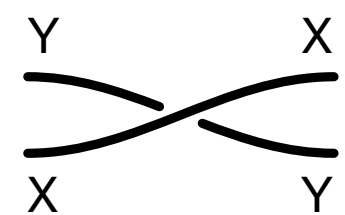


4. Symmetric Monoidal Categories

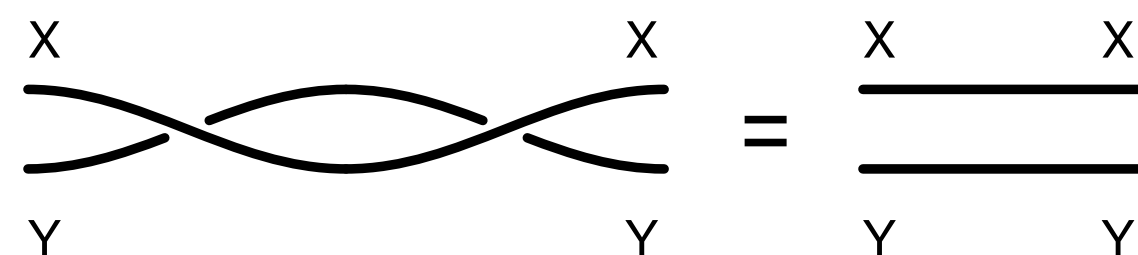
- we have a **swap** operation



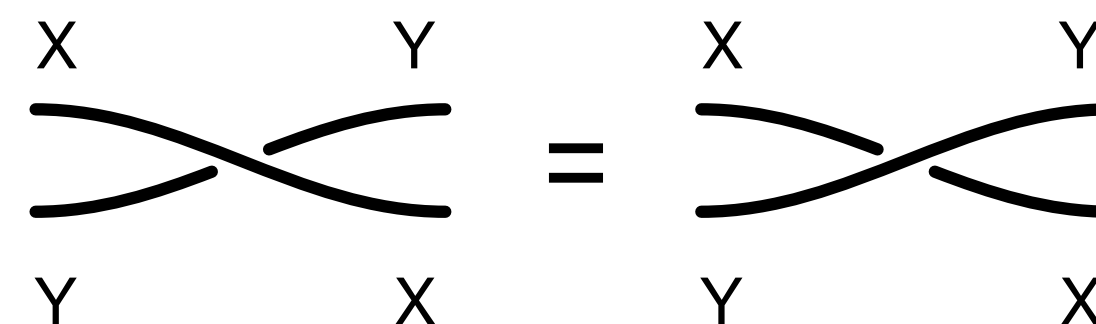
- with inverse depicted:



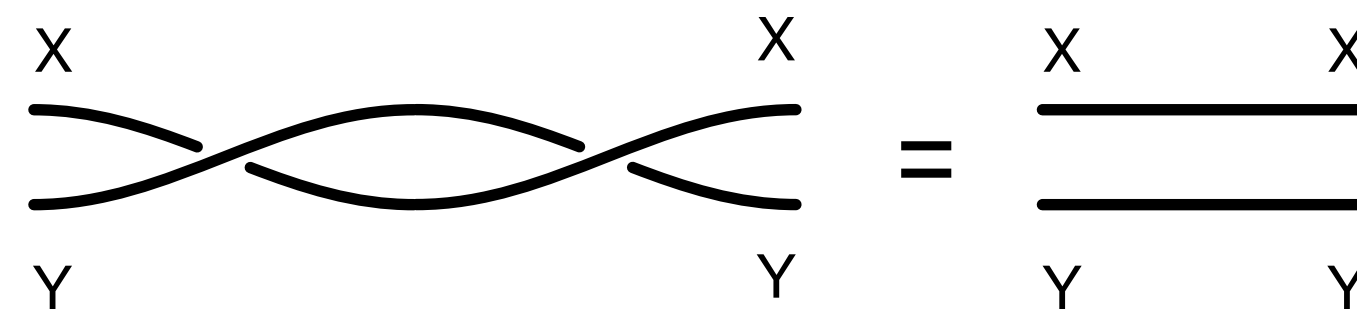
- this representation is chosen so that:



- furthermore, the swap satisfies



- i.e.



Circuit diagrams

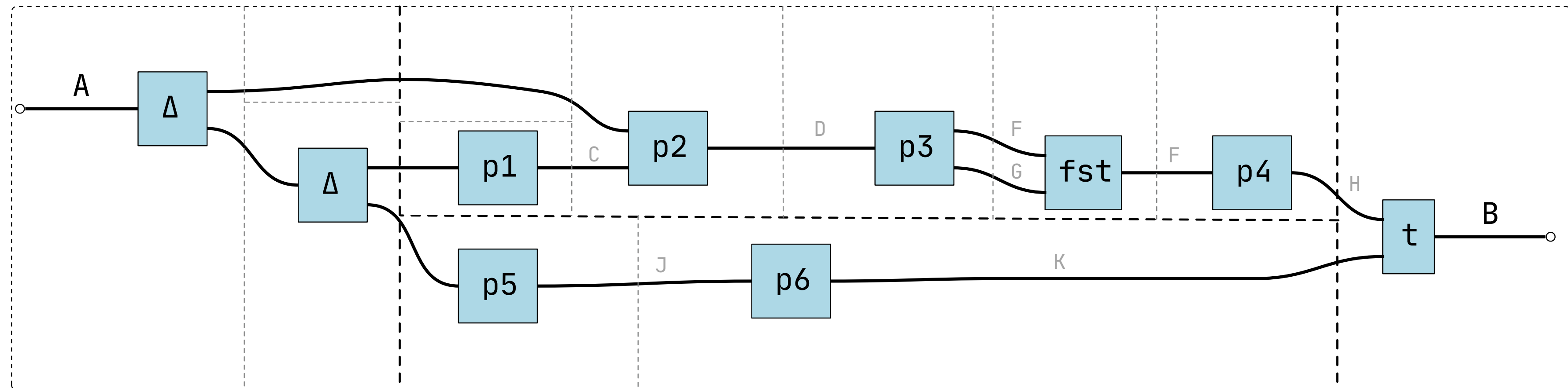
- The diagrams that we have discussed so far contain no loops.
- To be more precise they are called "circuit diagrams"
- Feedback loops can be introduced via **compact closed categories** (which won't be discussed here)

Part III: String Diagrams in Scala

Airflow-like process manager

```
val top = (id ++ p1) >>> p2 >>> p3 >>> fst >>> p4
```

```
val initial =  
  Δ >>> (id ++ Δ) >>> assocL
```



```
val bottom = p5 >>> p6
```

```
val program = initial >>> (top ++ bottom) >>> t
```


Tagless Process DSL

Category Structure

```

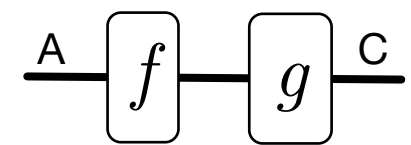
trait ProcessOpsDSL[Process[_], _] {

  type  $\rightsquigarrow$ [A, B] = Process[A, B]
  type  $\leftarrow$ [B, A] = Process[A, B]

  def id[A]: A  $\rightsquigarrow$  A

  extension [A, B, C] (f: A  $\rightsquigarrow$  B)
    @alpha("andThen")
    def >>> (g: B  $\rightsquigarrow$  C): A  $\rightsquigarrow$  C

  extension [A, B, C] (g: B  $\rightsquigarrow$  C)
    @alpha("compose")
    def  $\circ$  (f: A  $\rightsquigarrow$  B): A  $\rightsquigarrow$  C = f >>> g
  
```



Cartesian Structure

```

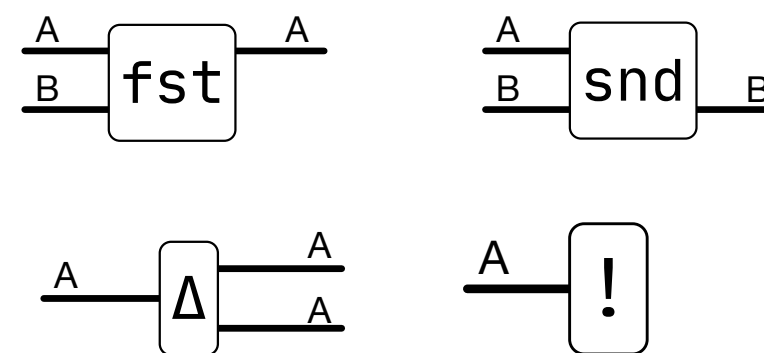
def fst[A, B]: (A, B)  $\rightsquigarrow$  A
def snd[A, B]: (A, B)  $\rightsquigarrow$  B

extension [A, B, C] (f: A  $\rightsquigarrow$  B)
  @alpha("mergeInput")
  def &&& (g: A  $\rightsquigarrow$  C): A  $\rightsquigarrow$  (B, C)

  @alpha("duplicate")
  def  $\Delta$ [A]: A  $\rightsquigarrow$  (A, A) =
    id[A] &&& id[A]

  // terminal object and monoidal unit
  type I = EmptyTuple

  def discard[A]: A  $\rightsquigarrow$  I
  
```



Monoidal Structure

```

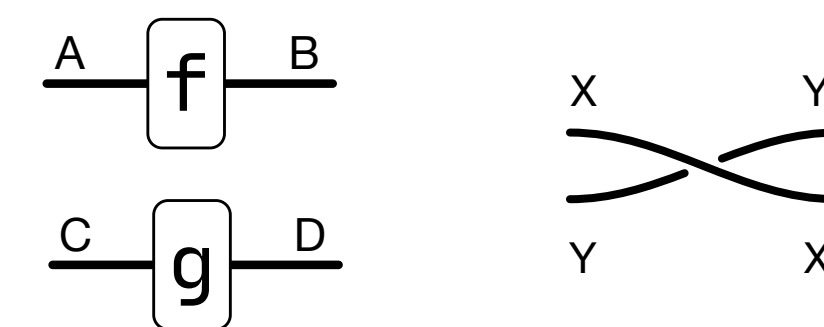
def assocR[X, Y, Z] : ((X, Y), Z)  $\rightsquigarrow$  (X, (Y, Z))
def assocL[X, Y, Z] : ((X, Y), Z)  $\leftarrow$  (X, (Y, Z))

def injR[X]: X  $\rightsquigarrow$  (I, X)
def injL[X]: X  $\rightsquigarrow$  (X, I)

// tensor on objects: A  $\otimes$  B = (A, B)
// tensor on arrows: f  $\otimes$  g = f ++ g

extension [A1, A2, B1, B2] (f: A1  $\rightsquigarrow$  B1)
  @alpha("combine")
  def ++ (g: A2  $\rightsquigarrow$  B2): (A1, A2)  $\rightsquigarrow$  (B1, B2)

def swap[X, Y] : (X, Y)  $\rightsquigarrow$  (Y, X)
def swapInverse[X, Y]: (X, Y)  $\leftarrow$  (Y, X)
  
```



Laws (the fine print)

Category Laws

```
@Law
def associativity[A, B, C, D](
  f: A  $\rightsquigarrow$  B,
  g: B  $\rightsquigarrow$  C,
  h: C  $\rightsquigarrow$  D
) =
  h  $\circ$  (g  $\circ$  f)  $\leftrightarrow$  (h  $\circ$  g)  $\circ$  f

@Law
def identityL[A, B](f: A  $\rightsquigarrow$  B) =
  (id[B]  $\circ$  f)  $\leftrightarrow$  f

@Law
def identityR[A, B](f: A  $\rightsquigarrow$  B) =
  f  $\circ$  id[A]  $\leftrightarrow$  f
```

Tensor Laws

```
@Law
def tensorCompositionLaw[
  A1, A2, B1, B2, C1, C2
](
  f1: A1  $\rightsquigarrow$  B1, f2: A2  $\rightsquigarrow$  B2,
  g1: B1  $\rightsquigarrow$  C1, g2: B2  $\rightsquigarrow$  C2,
) =
  (g1  $\circ$  f1) ++ (g2  $\circ$  f2)  $\leftrightarrow$ 
  (g1 ++ g2)  $\circ$  (f1 ++ f2)

@Law
def tensorIdentityLaw[A, B] =
  id[A] ++ id[B]  $\leftrightarrow$  id[(A, B)]
```

Monoidal Laws

```
@Law
def triangleEquation[X, Y] =
  assocR[X, I, Y] >>> (id[X] ++ snd[I, Y])  $\leftrightarrow$ 
  (fst[I, Y] ++ id[Y])

@Law
def pentagonEquation[W, X, Y, Z] =
  (assocR[W, X, Y] ++ id[Z]) >>>
  assocR[W, (X, Y), Z] >>>
  (id[W] ++ assocR[X, Y, Z])  $\leftrightarrow$ 
  (assocR[(W, X), Y, Z] >>> assocR[W, X, (Y, Z)])
```

Symmetric Monoidal Laws

```
@Law
def hexagonEquation1[X, Y, Z] =
  (assocL[X, Y, Z] >>> (swap[X, Y] ++ id[Z]) >>>
  assocR[Y, X, Z] >>> (id[Y] ++ swap[X, Z]) >>>
  assocL[Y, Z, X])  $\leftrightarrow$ 
  swap[X, (Y, Z)]

@Law
def hexagonEquation2[X, Y, Z] =
  (assocR[X, Y, Z] >>> (id[X] ++ swap[Y, Z]) >>>
  assocL[X, Z, Y] >>> (swap[X, Z] ++ id[Y]) >>>
  assocR[Z, X, Y])  $\leftrightarrow$ 
  swap[(X, Y), Z]

@Law
def symmetry[X, Y] =
  swap[X, Y]  $\leftrightarrow$ 
  swapInverse[Y, X]
```

ZIO interpreter

```
given RIOProcessOps as ProcessDSL0ps[RIO]:
```

```
def id[A] = RIO.identity[A]
```

```
@alpha("andThen")
```

```
def [A, B, C] (f: RIO[A, B]) >>> (g: RIO[B, C]) = f andThen g
```

```
def fst[A, B]: RIO[(A, B), A] = RIO.first
```

```
def snd[A, B]: RIO[(A, B), B] = RIO.second
```

```
@alpha("mergeInput")
```

```
def [A, B, C] (f: RIO[A, B]) &&& (g: RIO[A, C]): RIO[A, (B, C)] = f &&& g
```

```
def discard[A]: RIO[A, I] = RIO.succeed(EmptyTuple)
```

```
def assocR[X, Y, Z] : ((X, Y), Z) ↗ (X, (Y, Z)) = RIO.fromFunction { case ((x, y), z) => (x, (y, z)) }
```

```
def assocL[X, Y, Z] : ((X, Y), Z) ↖ (X, (Y, Z)) = RIO.fromFunction { case (x, (y, z)) => ((x, y), z) }
```

```
def injR[X]: X ↗ (I, X) = RIO.fromFunction { a => (EmptyTuple, a) }
```

```
def injL[X]: X ↖ (X, I) = RIO.fromFunction { a => (a, EmptyTuple) }
```

```
@alpha("combine")
```

```
def [A1, A2, B1, B2] (f: A1 ↗ B1) ++ (g: A2 ↗ B2): (A1, A2) ↗ (B1, B2) =
```

```
  val left  = fst >>> f
```

```
  val right = snd >>> g
```

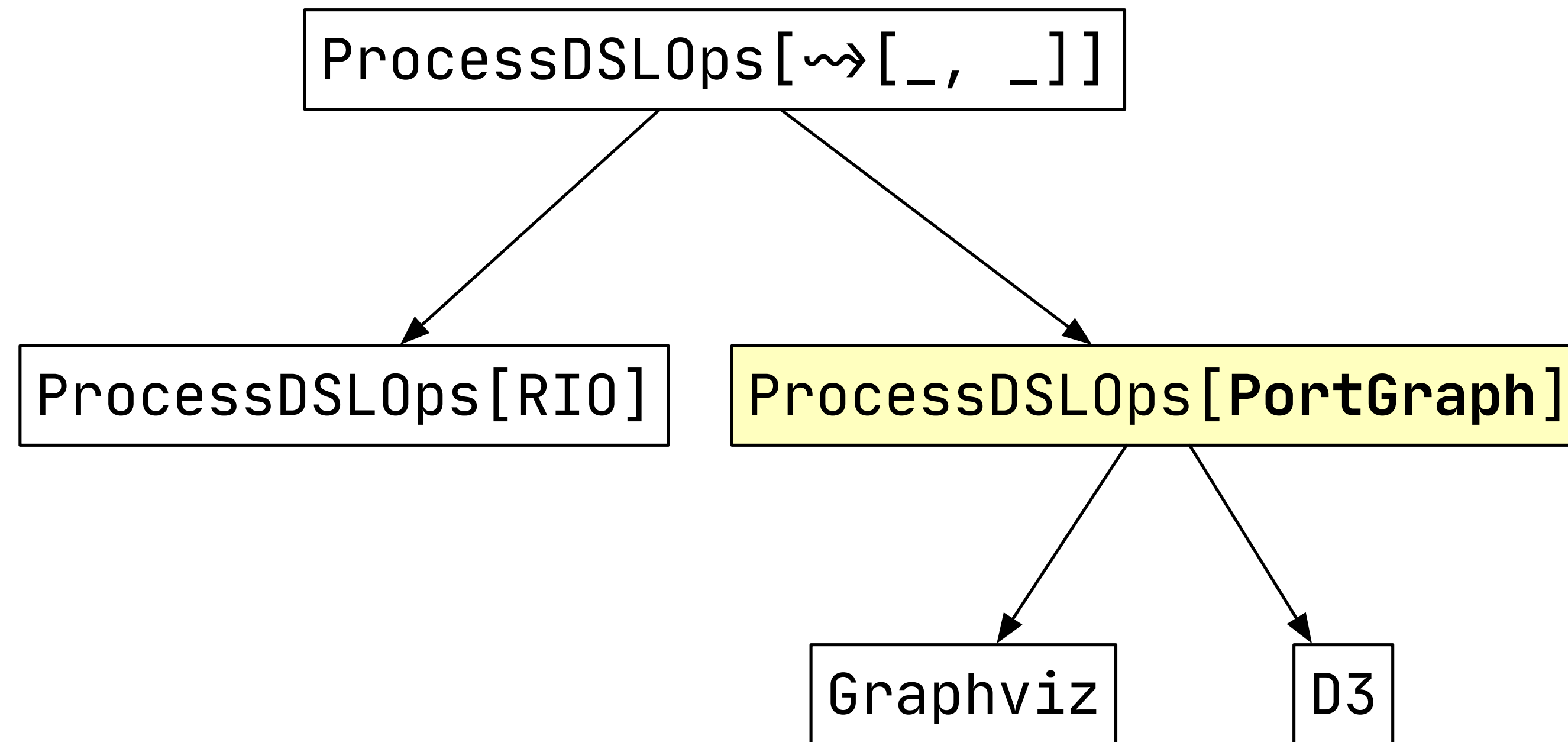
```
  (left zipWithPar right) ((x, y) => (x, y))
```

```
def swap [X, Y]: (X, Y) ↗ (Y, X) = RIO.swap
```

```
def swapInverse[X, Y]: (X, Y) ↖ (Y, X) = RIO.swap
```

RIO[A, B]
≈
A ⇒ Either[Throwable, B]

Process DSL



Category of Port Graphs

```

val graph1 =
  PortGraph(
    boxes = SeqMap(
      a → (List("A"), List("B", "C", "D")),
      b → (List("D", "C", "G"), List("Y", "E", "F")),
      c → (List("B", "Y"), List("Z")),
    ),
    incoming = List((a, 0), (b, 2)),
    inner = List(
      (a, 0) → (c, 0),
      (a, 1) → (b, 1),
      (a, 2) → (b, 0),
      (b, 0) → (c, 1),
    ),
    outgoing = List((c, 0), (b, 1), (b, 2))
  )

```

```

val graph2 = singleBox("x", List("Z", "E", "F"), List("W"))

```

```

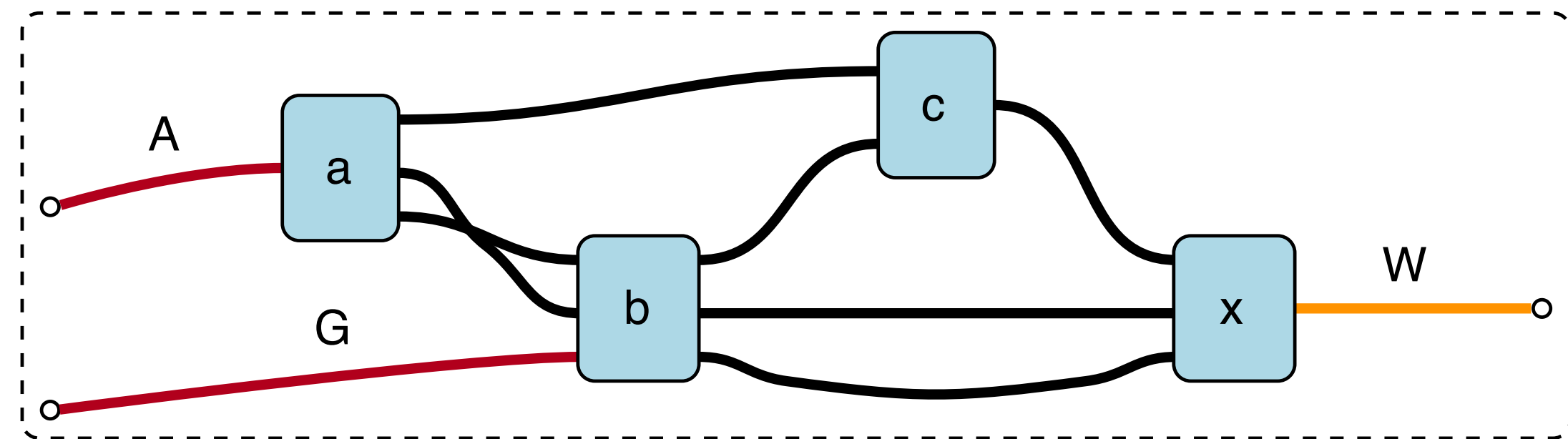
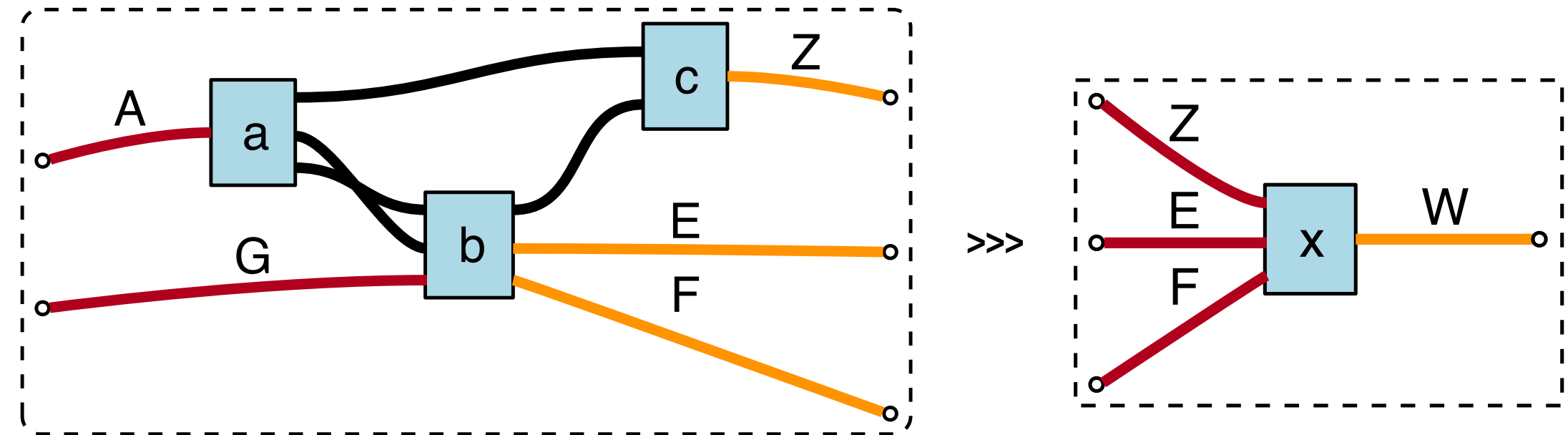
val combined = graph1 andThen graph2

```

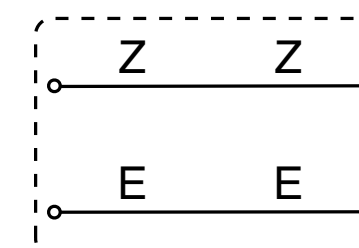
```

def render[B, P](pg: PortGraph[B, P]): graphviz.model.Graph
  = ???

```



Identity:



the end

Bibliography

- Brendan Fong David I. Spivak . Seven Sketches in Compositionality
- John C. Baez, Mike Stay. Physics, Topology, Logic and Computation: A Rosetta Stone
- Bob Coecke, Aleks Kissinger. Picturing Quantum Processes: A First Course in Quantum Theory and Diagrammatic Reasoning.
- Aleks Kissinger. Pictures of Processes

Extras

Tagless Process DSL

Constructor

```
trait ProcessDSL[↔[_, _]]:  
  def httpCall[A: BodySerializer, B: Decoder](  
    name: String, uri: Uri  
  ): A ↔ B
```

ZIO instance

```
given RIOProcessDSL as ProcessDSL[RIO]:  
  val backend = HttpURLConnectionBackend()  
  import sttp.client3._  
  import sttp.client3.circe._  
  
  def httpCall[A: BodySerializer, B: Decoder](  
    name: String, uri: Uri  
  ): RIO[A, B] =  
    RIO.fromFunctionM[A, B] { (a: A) =>  
      RIO.fromEither {  
        basicRequest  
          .post(uri)  
          .body(a)  
          .response(asJson[B])  
          .send(backend)  
          .body  
          .left.map(new Exception(_))  
      }  
    }
```